

GRAPHICAL NOTATION SCHEMES

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WHAT'S THE ISSUE?

- Symbolic expressions used in foundational mathematics
 - Powerful methods
 - Objects of study in themselves
 - Can be technical and syntax-heavy
 - Can be easy to make mistakes by hand and hard to spot structure

WHAT'S THE ISSUE?

- Researchers have always found ways round this:
 - Doodles in margins to help symbolic calculations
 - Proofs-by-picture

WHAT'S THE ISSUE?

- Pictures can capture important aspects of abstract structure
 - It's not a coincidence that they're so useful
 - Many classes of graphs exhibit rich categorical structure
 - That's *why* they're useful!

STRING DIAGRAMS

- Just one type of example: monoidal categories with other structure
 - Nice examples to demonstrate the ideas

MONOIDAL CATEGORIES

- A **monoidal category** is a category with
 - a bifunctorial **tensor product**, \otimes
 - A specified “tensor unit”, I
 - Associativity and identity natural isomorphisms, a , l , r , or strictness
 - Coherence axioms

MONOIDAL CATEGORIES

- Examples
 - $(\mathbf{Set}, \times, \{*\})$ non-strict
 - (\mathcal{C}, \times, t) category with binary products and terminal object
 - \mathbf{Rel} category of sets and relations
 - $([\mathcal{C}, \mathcal{C}], \circ, \text{id}_{\mathcal{C}})$ strict

MONOIDAL EXPRESSIONS

Expressions in categories

$$f : A \rightarrow B \quad g : B \rightarrow C \quad g \circ f : A \rightarrow C$$

$$\text{“}C\text{”} = \text{id}_C : C \rightarrow C$$

Expressions in monoidal categories

$$f \otimes g : A \otimes C \rightarrow B \otimes C \quad f \otimes C : A \otimes C \rightarrow B \otimes C$$

$$h : A \otimes B \rightarrow C \otimes D \otimes E \otimes F$$

$$f : X \rightarrow X_1 \otimes \cdots \otimes X_n$$

$$g : X_1 \otimes \cdots \otimes X_n \rightarrow Y$$

$$g \circ f : X \rightarrow Y$$

MONOIDAL EXPRESSIONS

- When are these equal *due to monoidal axioms*?
- When are these equal *in any monoidal category*?

$$f : A \rightarrow E \otimes D$$

$$h : D \rightarrow G \otimes H$$

$$g : D \otimes B \otimes C \rightarrow F$$

- (Again, working strictly)


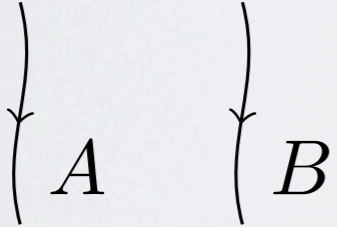
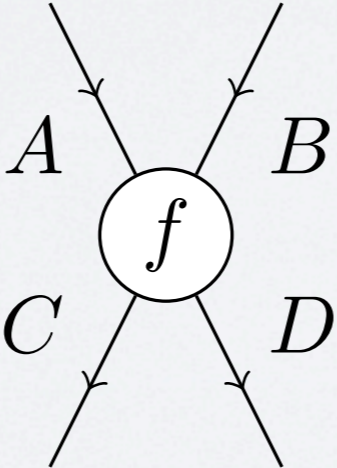
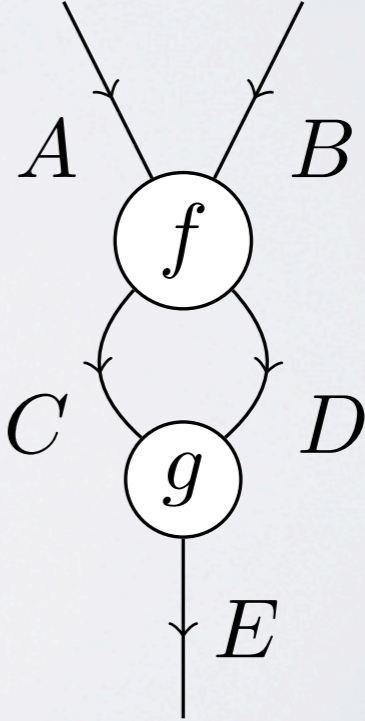
$$(E \otimes g \otimes h) \circ (f \otimes B \otimes C \otimes D)$$

$\parallel ?$

$$(E \otimes g \otimes G \otimes H) \circ (f \otimes B \otimes C \otimes h)$$

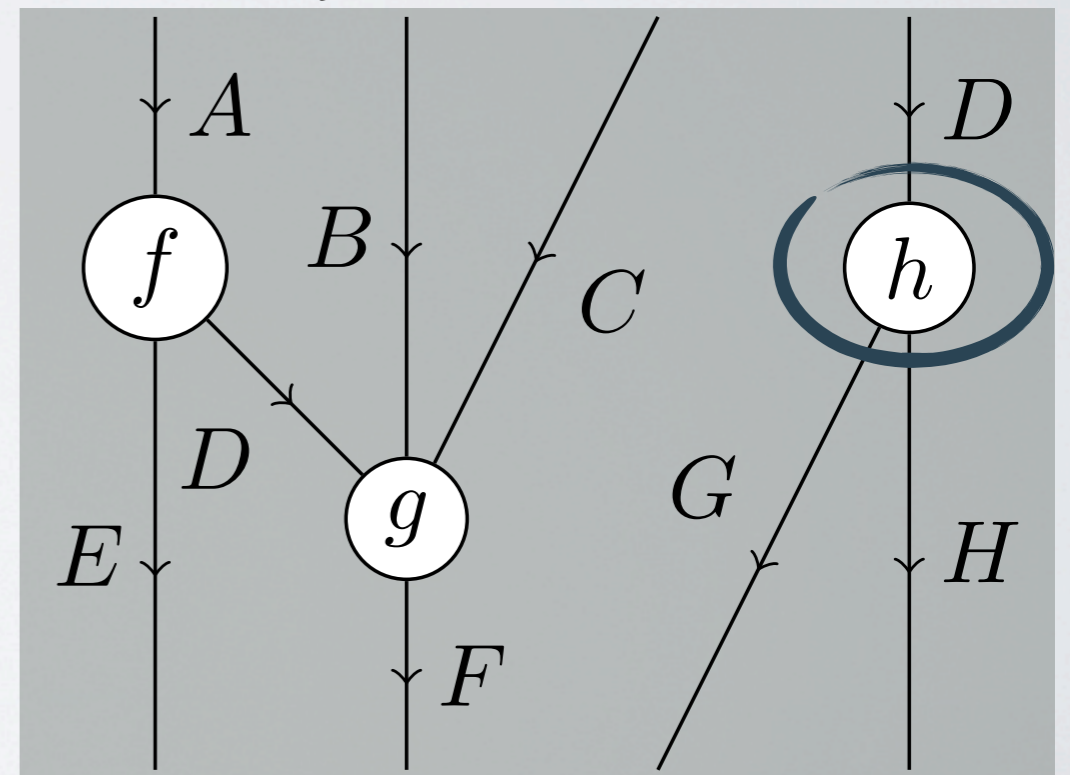
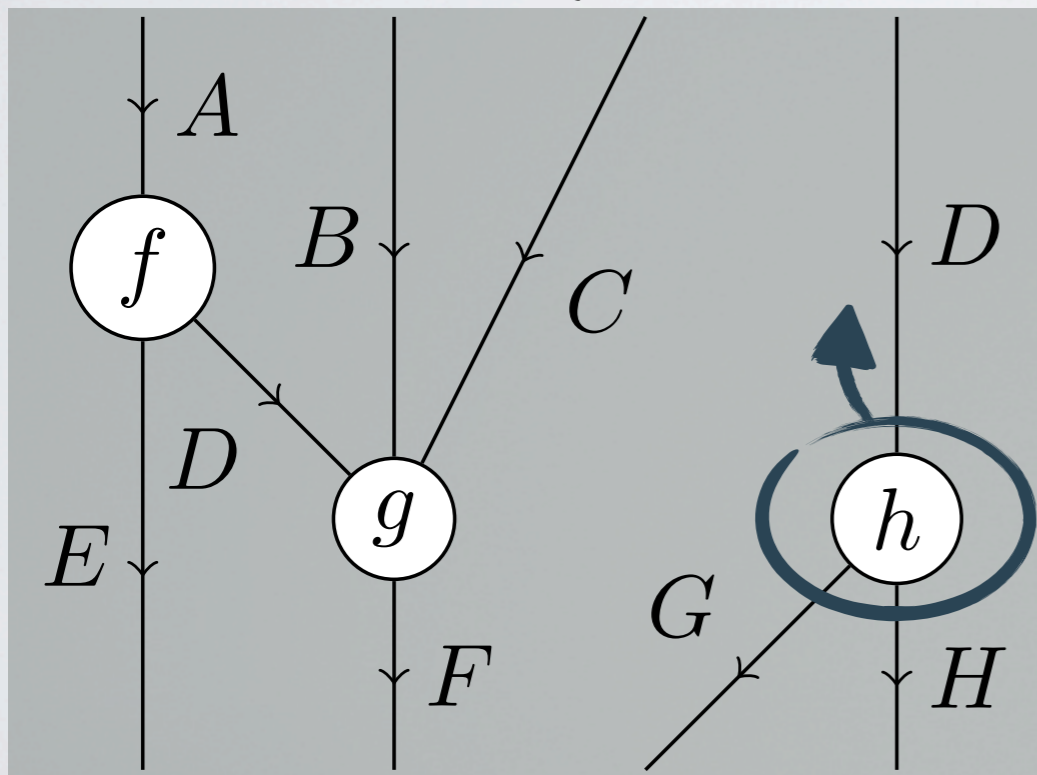
STRING DIAGRAMS FOR MONOIDAL CATEGORIES

- Diagrams to represent expressions in monoidal categories

| A | $A \otimes B$ | $A \otimes B \xrightarrow{f} C \otimes D$ | $A \otimes B \xrightarrow{f} C \otimes D \xrightarrow{g} E$ |
|---|---|---|---|
|  |  |  |  |

STRING DIAGRAMS FOR MONOIDAL CATEGORIES

$$(E \otimes g \otimes h) \circ (f \otimes B \otimes C \otimes D) = (E \otimes g \otimes G \otimes H) \circ (f \otimes B \otimes C \otimes h)$$



Full treatment: André Joyal and Ross Street's *Geometry of tensor calculus I*. Advances in Mathematics 1991

Hard to find online! :(

IDEA OF THEOREM

- **Theorem.** These symbolic expressions form a (free strict) monoidal category.
- **Theorem.** (Suitably-defined) labelled diagrams form a strict monoidal category.
- **Theorem.** These categories are monoidally equivalent.
 - Notion of diagram *valuation* and *evaluation*
 - Canonical diagram construction

STRING DIAGRAMS FOR MONOIDAL CATEGORIES

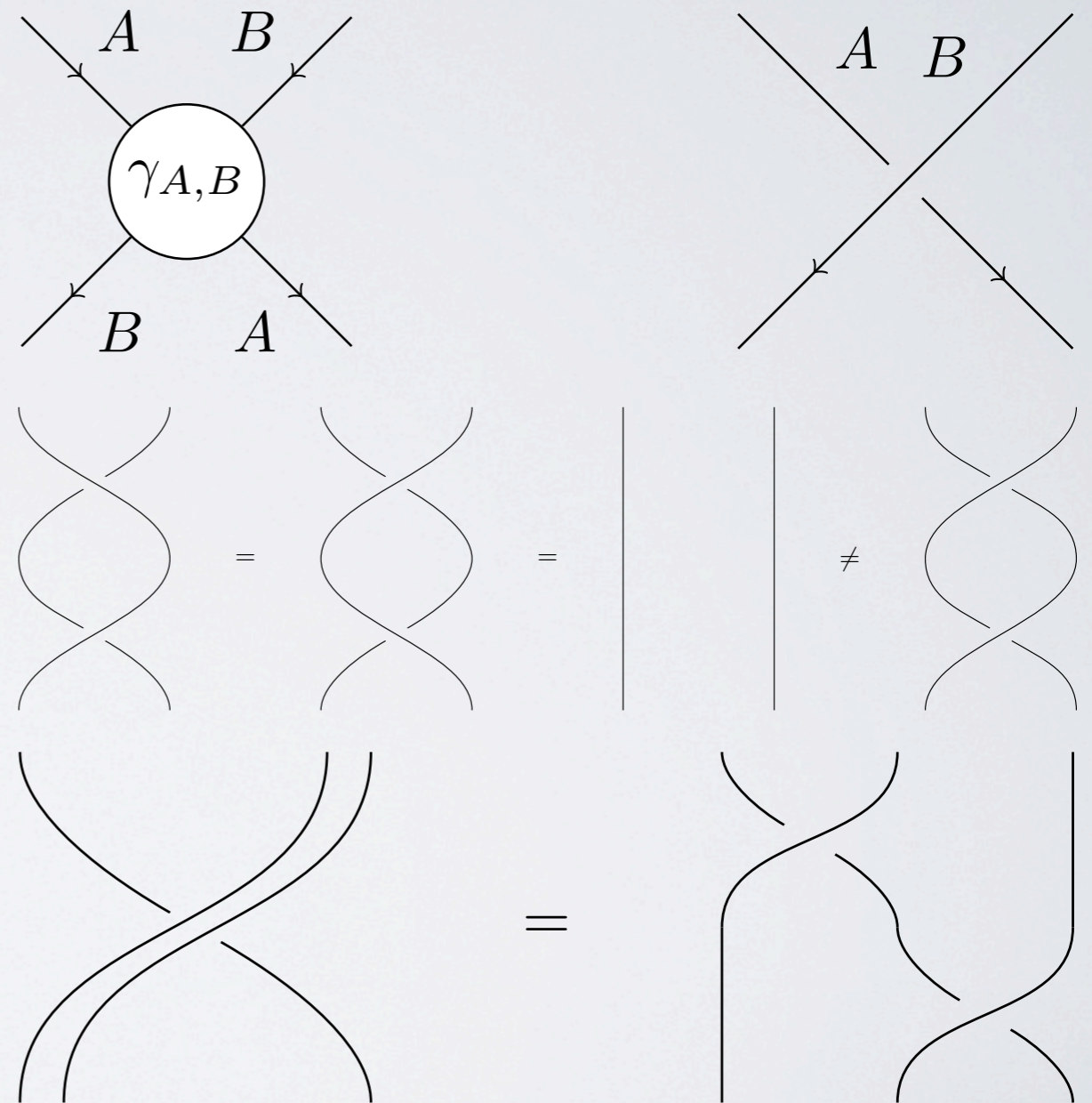
- This gives us:
 - Diagrams are a valid notation
 - Deformations on a diagram preserves valuation in category
 - We can do mathematics using these diagrams

ADDING STRUCTURE, AUGMENTING GRAPHS

- Can add more structure to a monoidal category
 - Can augment graphical language to capture new axioms
 - Some examples...

BRAIDING

- Add a **braiding** natural isomorphism
- Coherence
- Eg. category of braids



Full treatment: André Joyal and Ross Street's *Braided tensor categories*. Advances in Mathematics 1993

BRAIDING

- **Theorem.** (Joyal and Street) Free braided monoidal category is equivalent to category of labelled braids.
- **Theorem.** (Reidemeister) Manipulating braid diagrams corresponds exactly to isotopy on braided strings in 3-space.
- **Corollary.** (Joyal and Street) Two expressions are isomorphic iff the underlying braids are the same.

BRAIDING

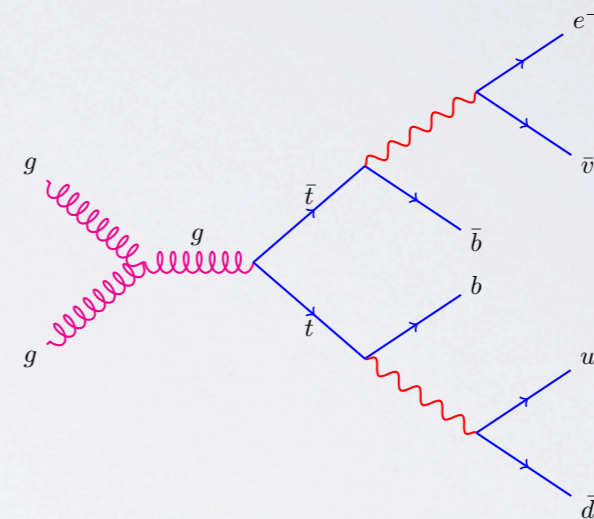
- We see some horrendous braiding isomorphisms...



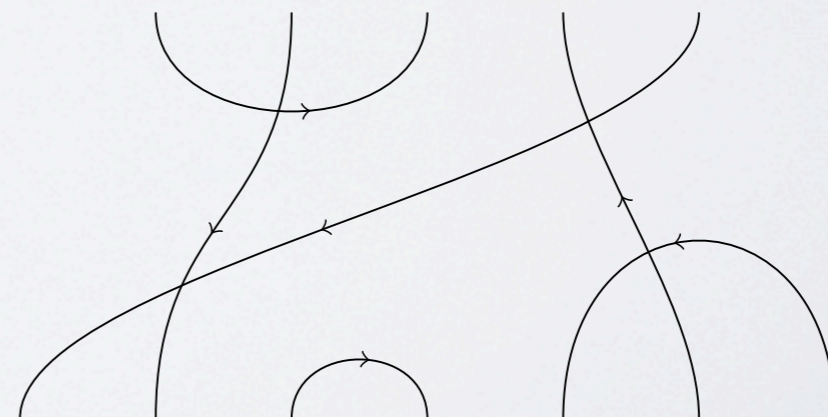
- ...are just the identity! We've saved a lot of chalk.

COMPACT CLOSED

- A **symmetry** is a self-inverse braiding
 - Sets; Feynman diagrams; ...
- A **compact closed** category is symmetric with (right) duals
 - Vector spaces; ...



dream.inf.ed.ac.uk/projects/quantomatic/talks/cambridge-2010-2x2.pdf



RIBBON

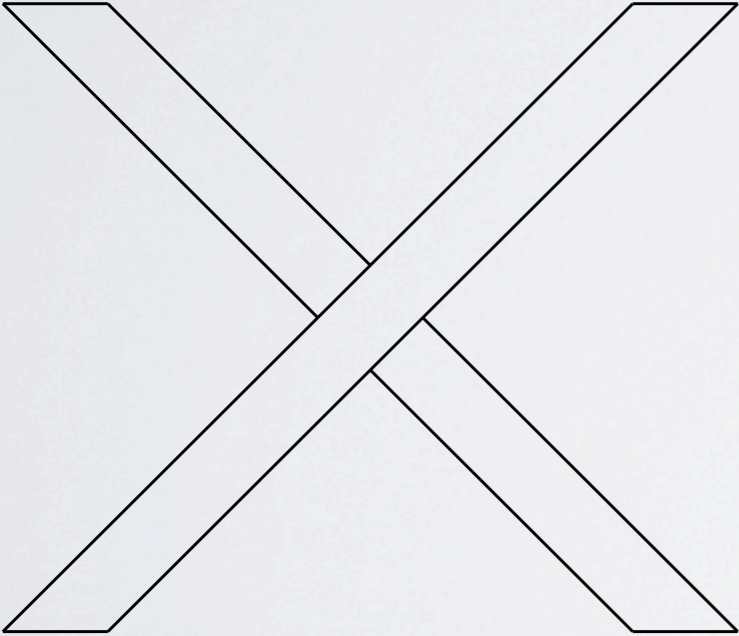
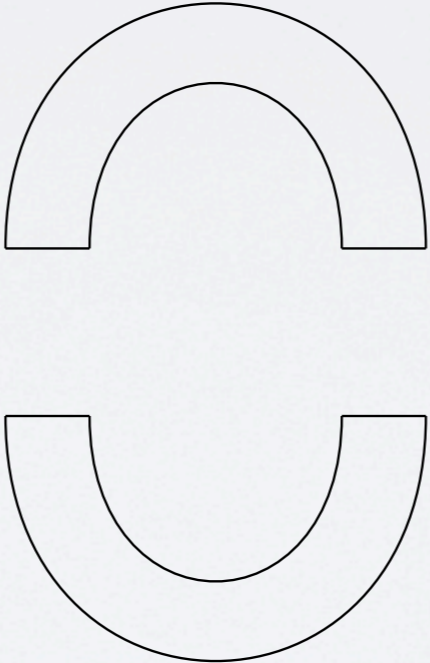
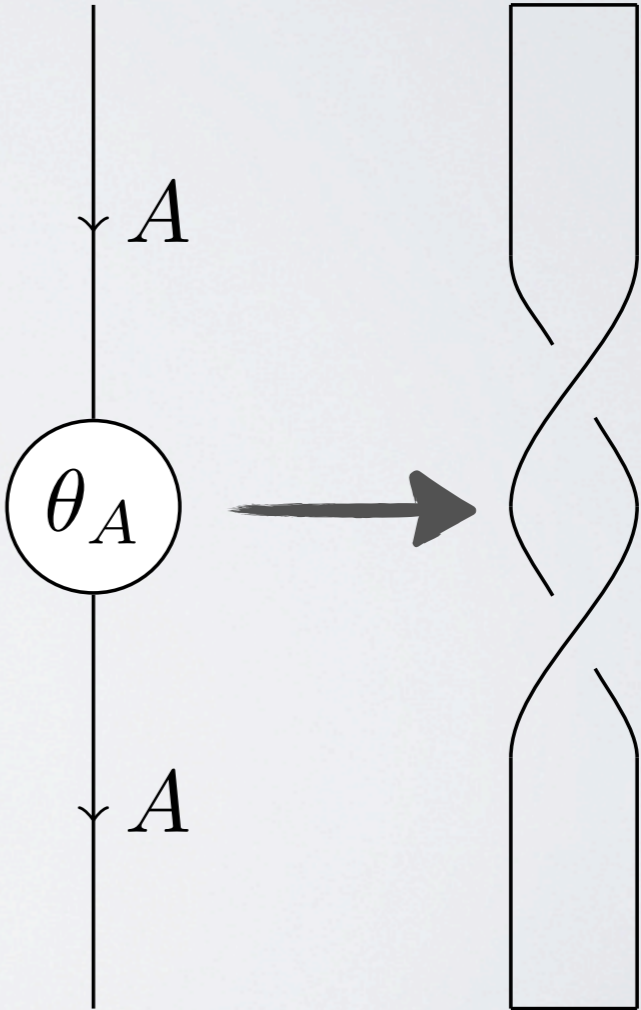
- A **twist** in a braided monoidal category is a natural isomorphism

$$\theta_A : A \rightarrow A$$

coherent with the braiding.

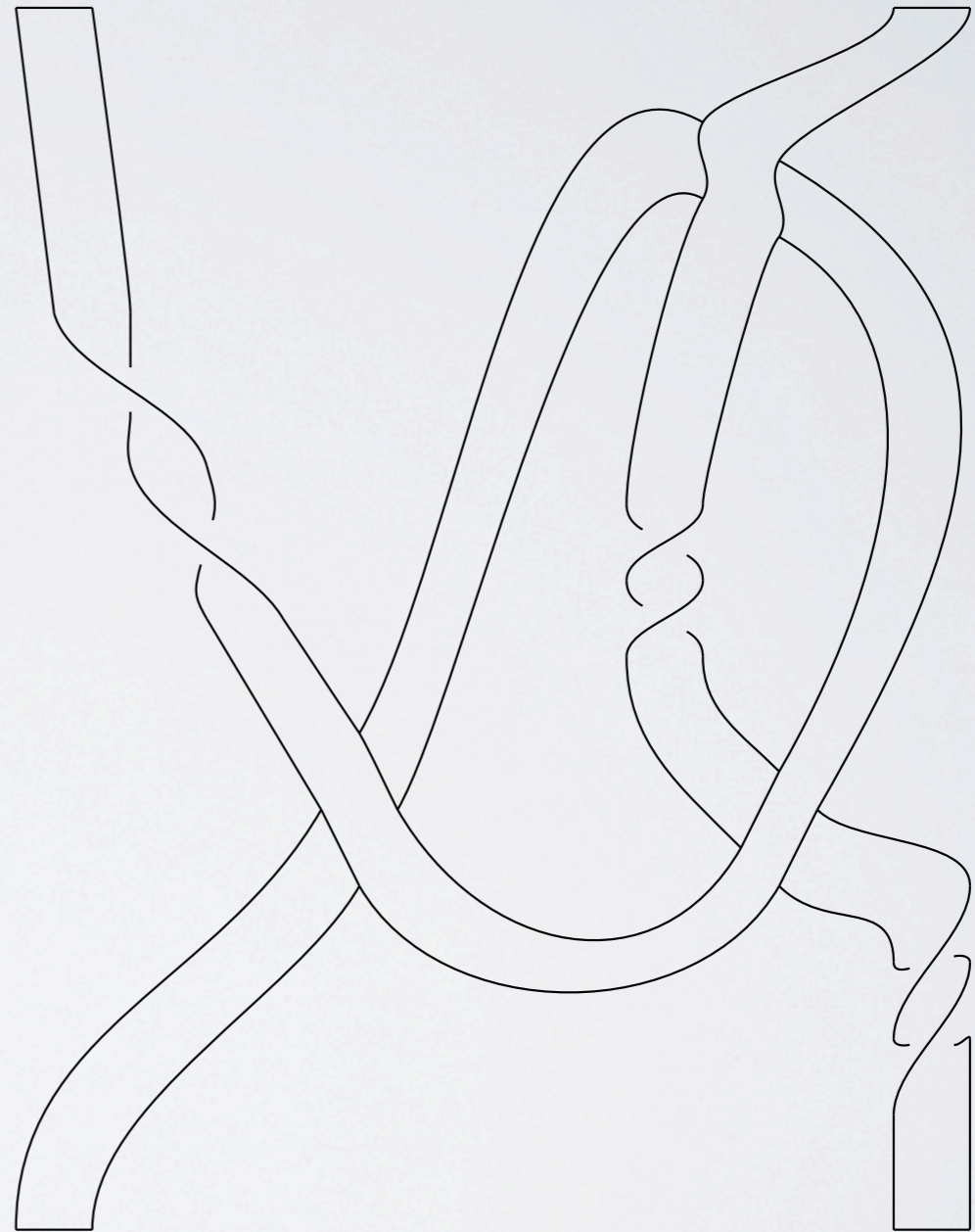
- A **tortile** or **ribbon category** is a braided monoidal category with a dual for each object and a twist (plus axioms)

RIBBON

| Braiding | Duals | Twist |
|---|---|--|
|  |  |  |

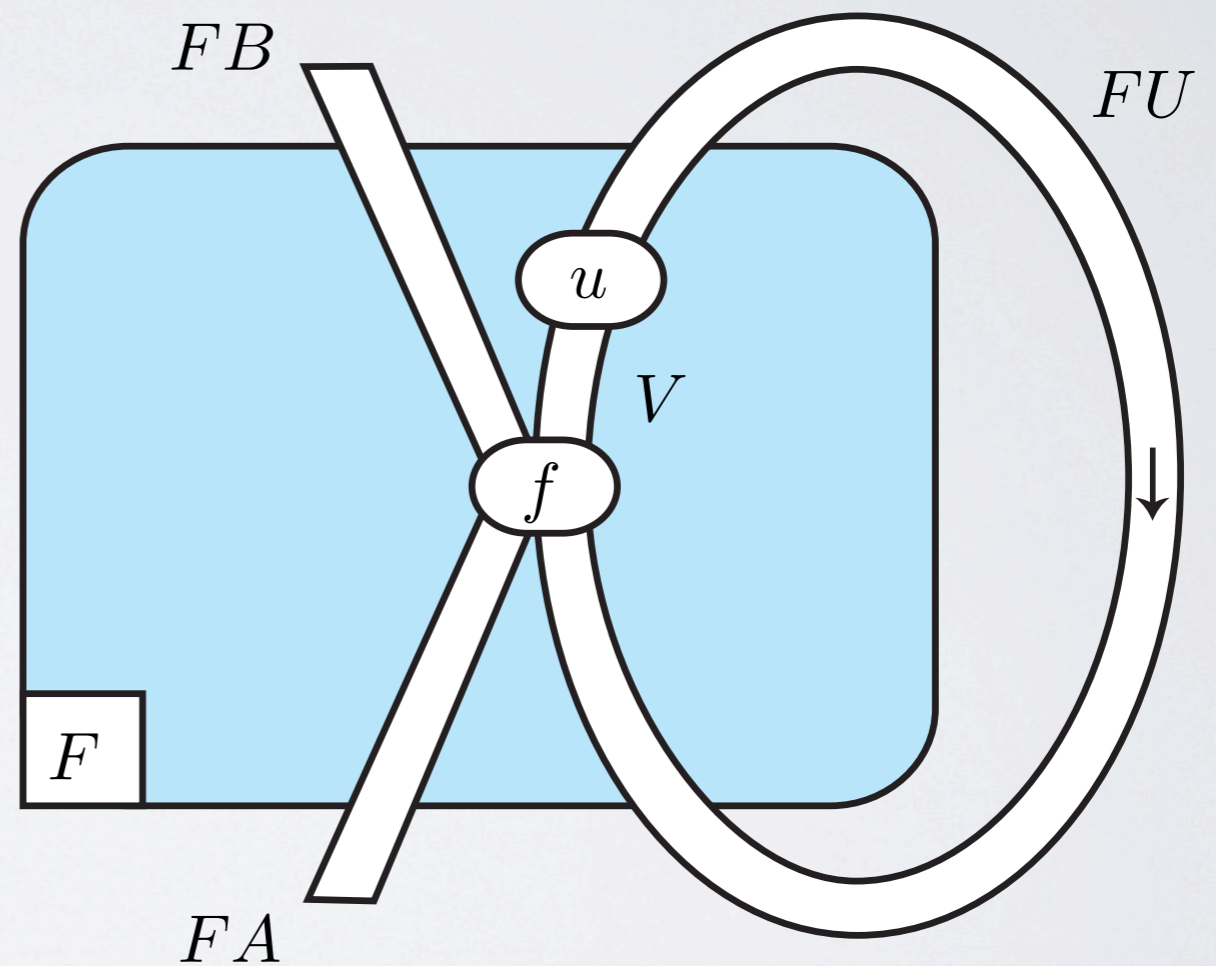
RIBBON

- These compose to form pictures
- Like ribbon tangles in 3-space!
 - (Missing a lot of detail again...)
- Useful for knot invariants, quantum protocols



EVEN FURTHER

- Functorial boxes
- Higher categories



WHAT I'M DOING

- Similar motivations:
 - Formalise graphical language people already use for argument
 - Not a monoidal category!

GAME SEMANTICS

- Model computational environment as interaction “games”
 - *Player* or *proponent* is system, *opponent* is environment
 - *Game* is alternating sequence of moves
- Games model types
- *Strategies* for player model terms

Many introductions around. Here's slides from a recent talk by Guy McCusker at LI2012:
li2012.univ-mrs.fr/media/talk19/mccusker-lectures.pdf

GAME SEMANTICS

- Example:

N

q

3

O

P

GAME SEMANTICS

- Example:

| | | | |
|-----|---------------|-----|-----|
| N | \rightarrow | N | |
| | | q | O |
| q | | | P |
| 3 | | | O |
| | | 4 | P |

GAME SEMANTICS

- Arrow games, $A \multimap B$
 - Two games in parallel
 - Roles reversed roles on left
 - Moves interleaved
 - Interleaving is a *schedule*

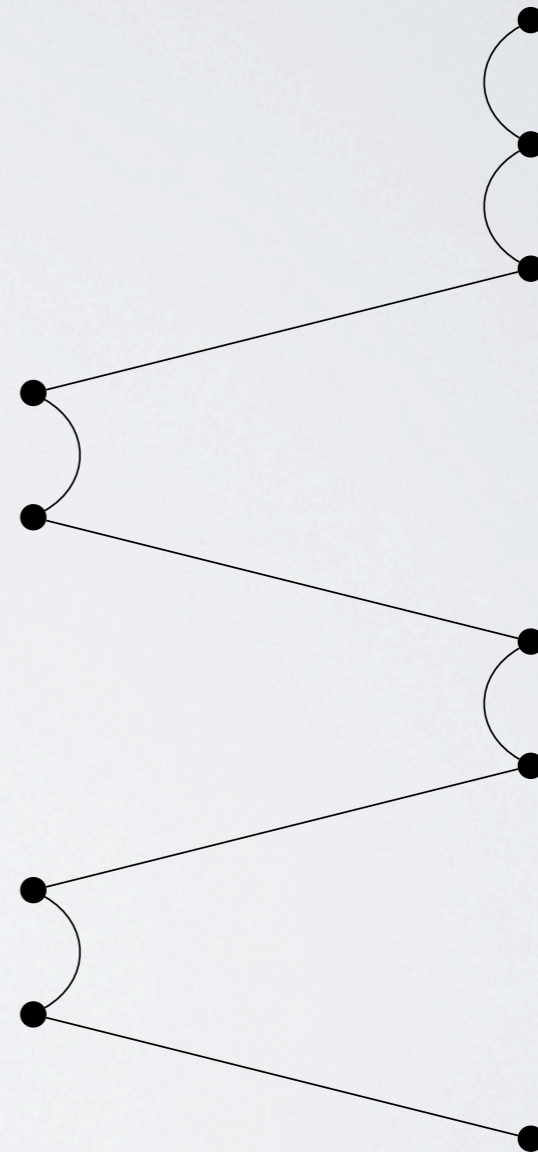
SCHEDULES

- Originally definition combinatorial in nature
 - Can be thought of as binary strings
 - Or “collectively surjective” function pairs, or order relations
- Composition is highly combinatorial (and tricky)
- Associativity is difficult to establish

Good stuff here: Russ Harmer, Paul-André Melliès and Martin Hyland's *Categorical combinatorics for innocent strategies*. LICS 2007
pps.jussieu.fr/~mellies/papers/lics2007-categorical-combinatorics.pdf

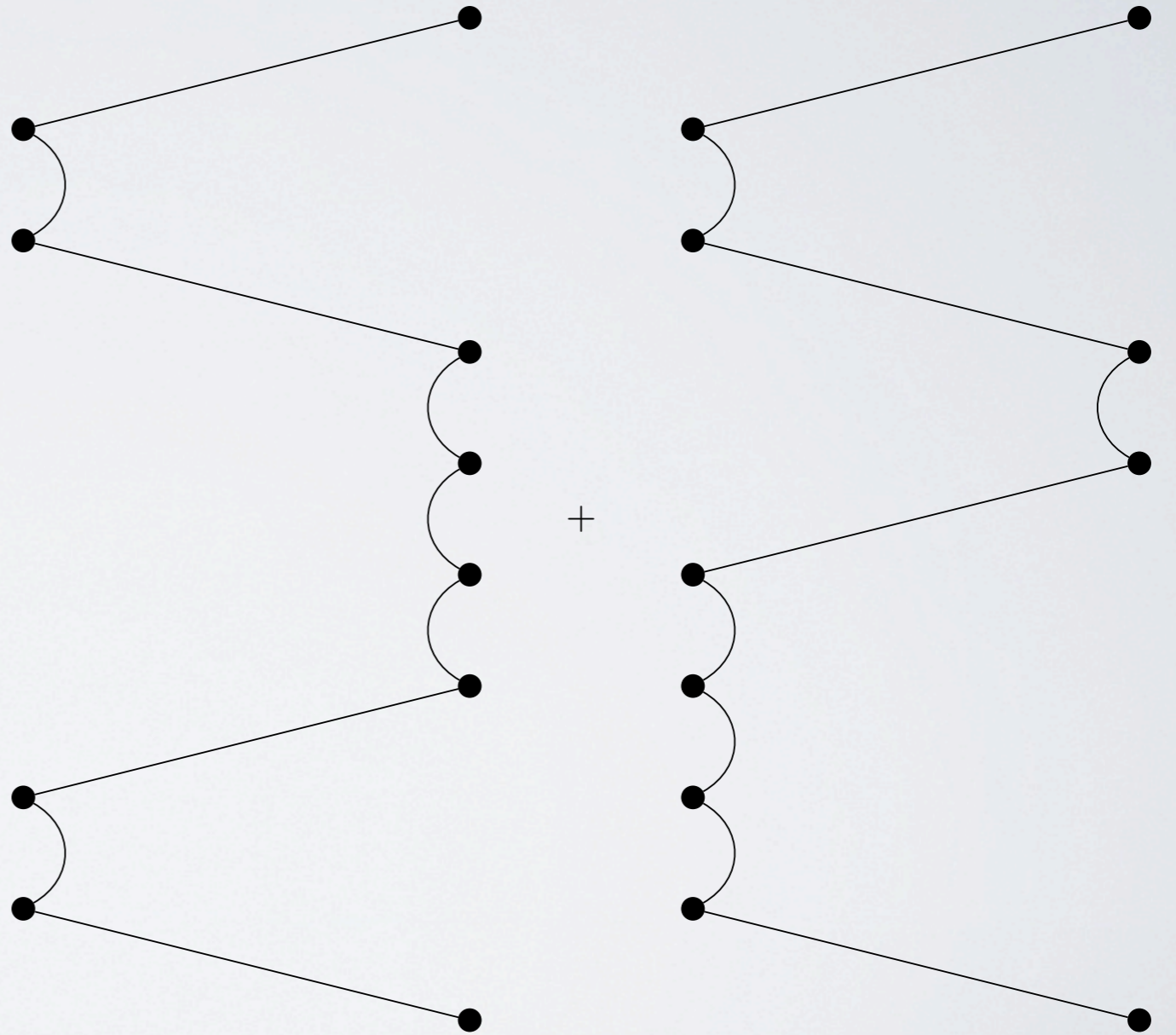
SCHEDULES

- Composition and associativity are tricky to do by hand
- People tend to use pictures



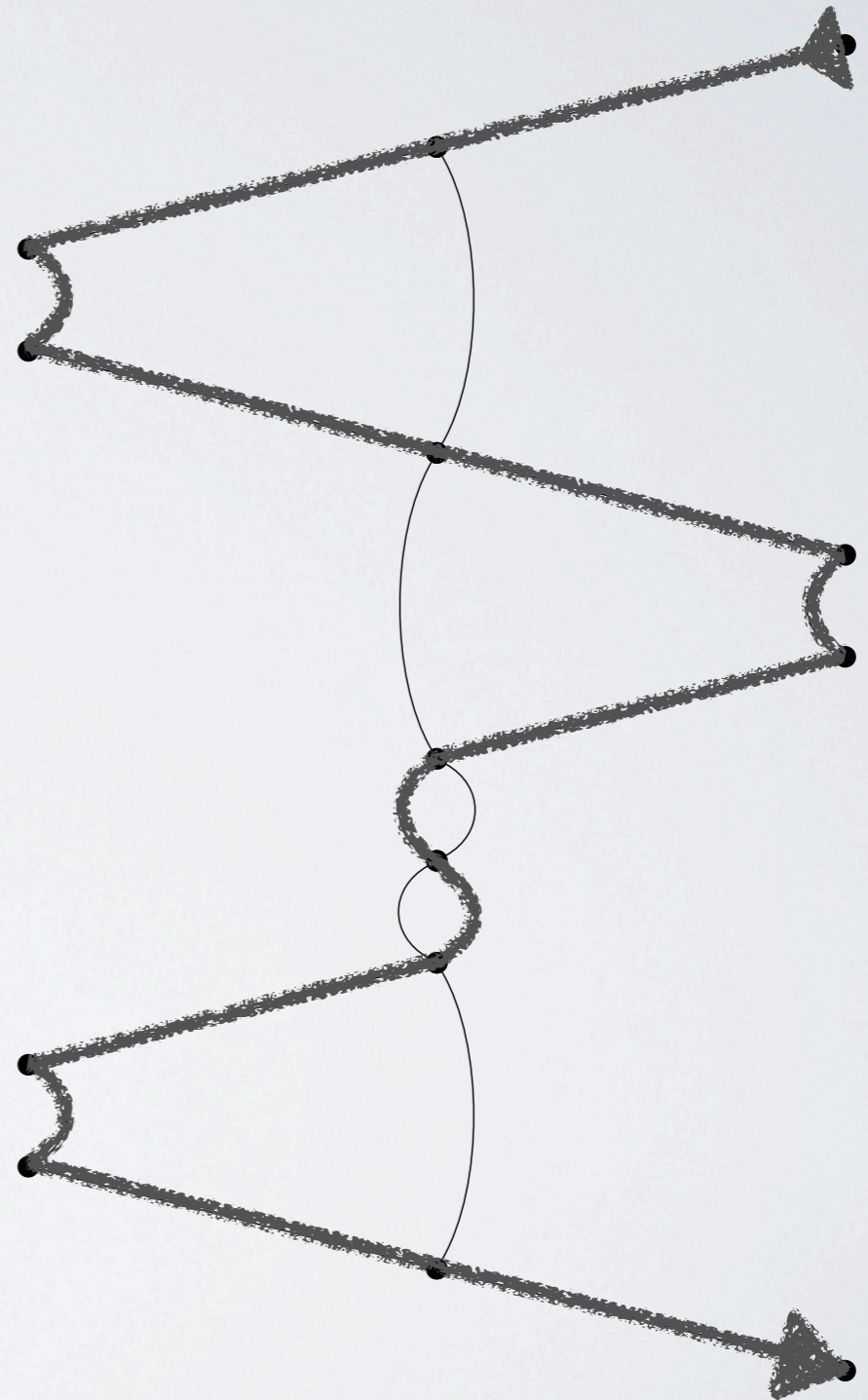
SCHEDULES

- Composition:



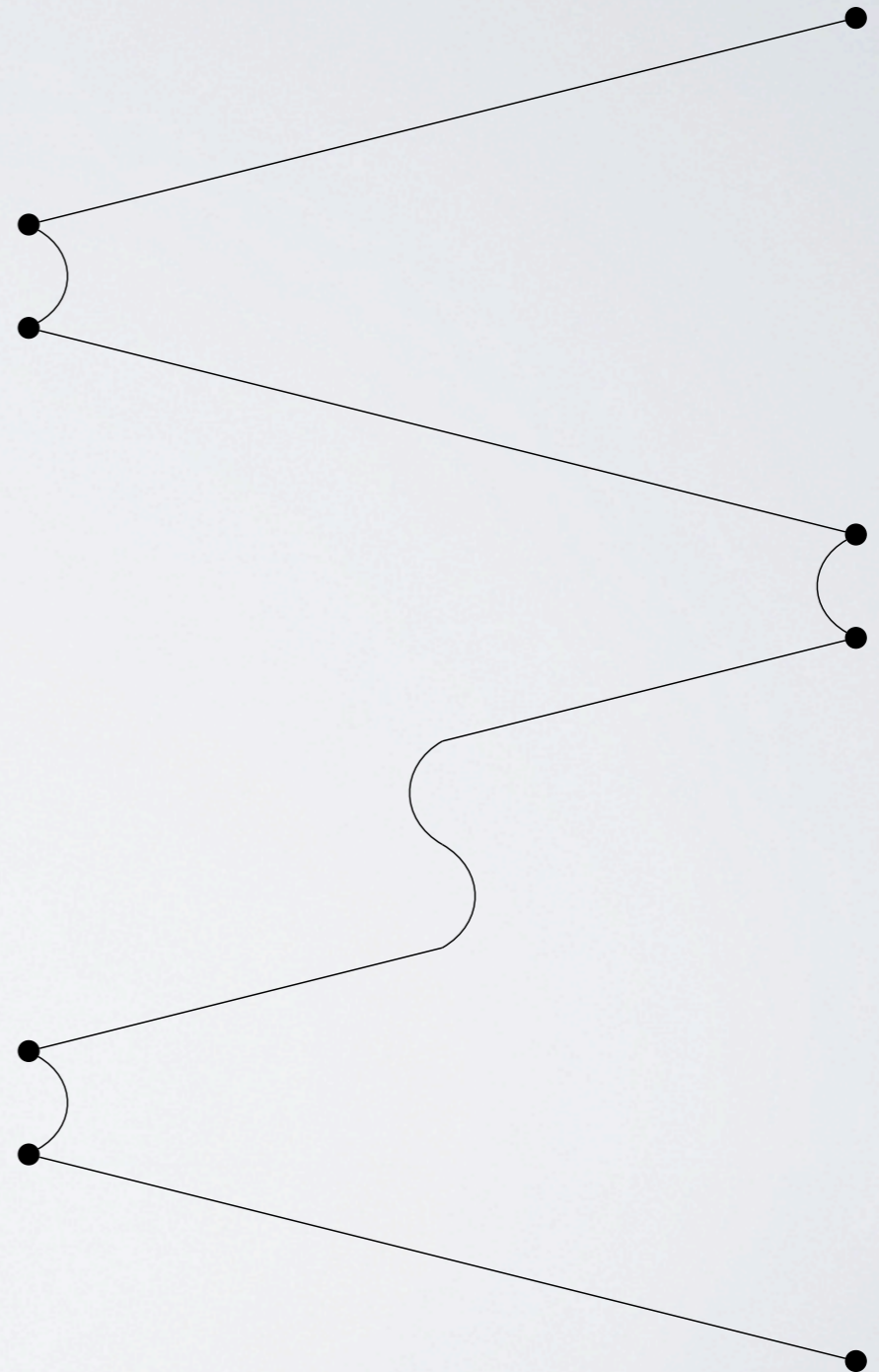
SCHEDULES

- Composition:
 - Glue schedules
 - Trace path through all nodes



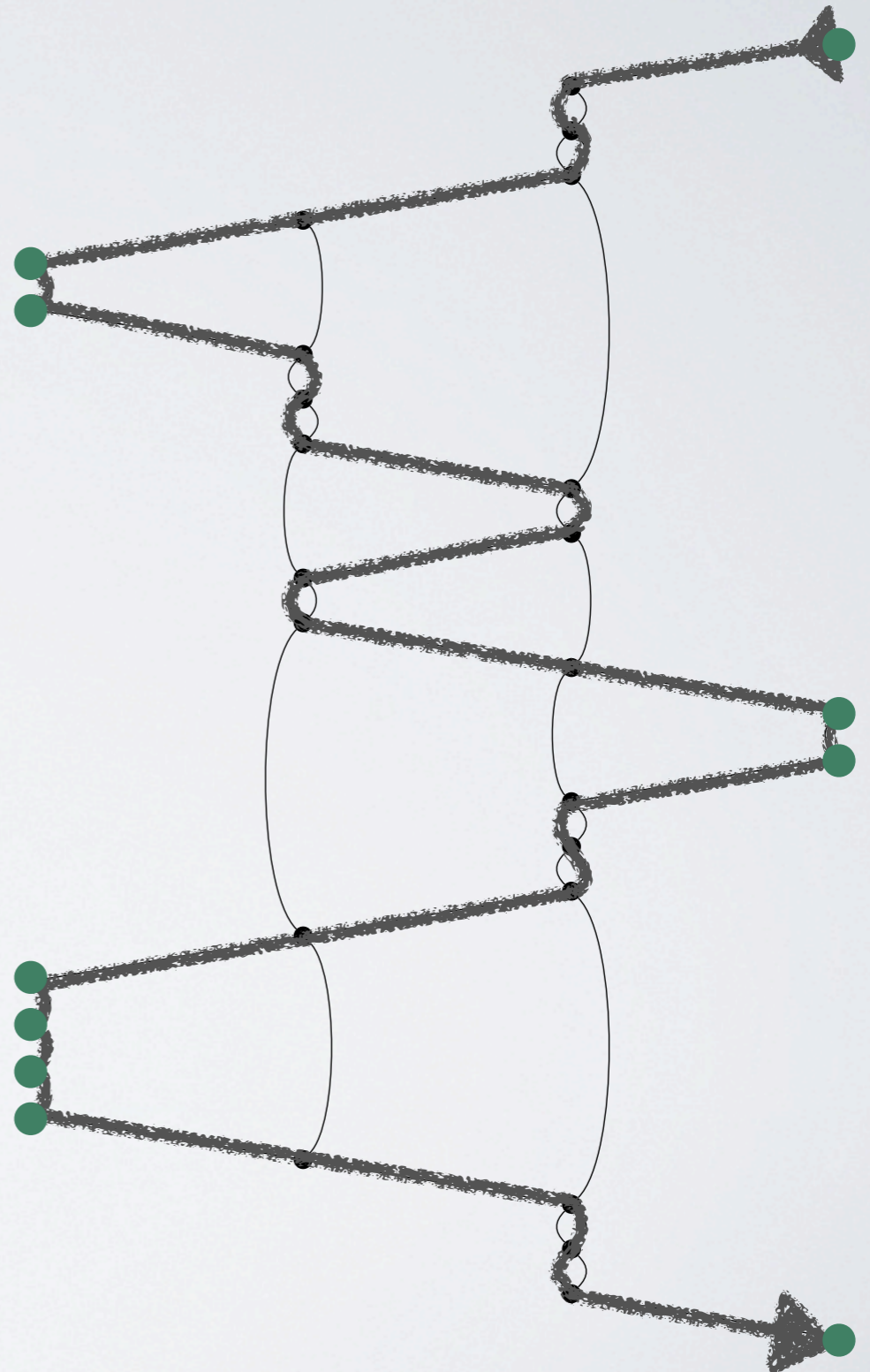
SCHEDULES

- Composition:
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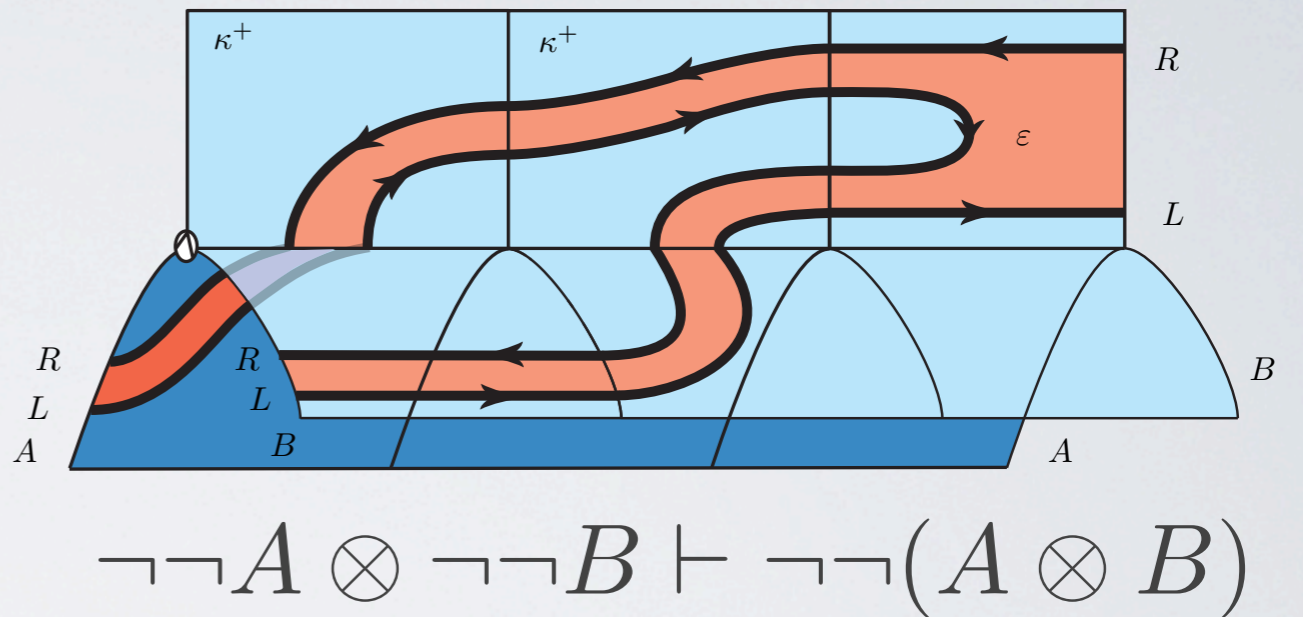
SCHEDULES

- Associativity becomes easy!
- “Juxtaposition in the plane is associative”



USES, RESEARCH

- Things I heard about at Logic and Interaction 2012

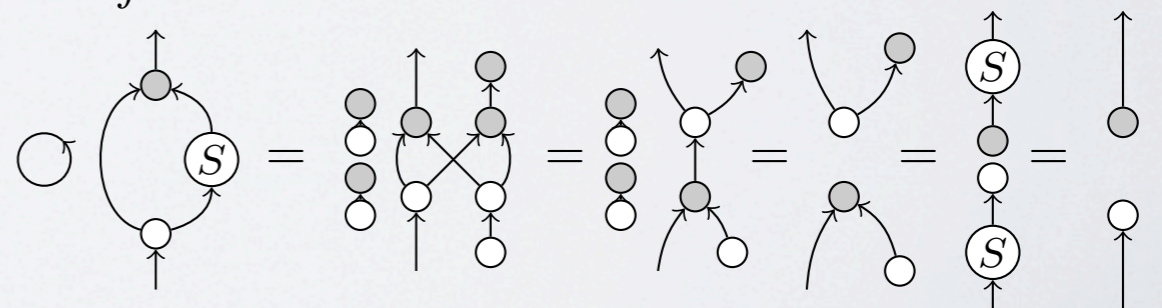


pps.jussieu.fr/~mellies/tensorial-logic.html

- Melliès' *Tensorial logic*
- Coecke, Duncan, Kissinger and Wang: categorical quantum mechanics

Theorem 3.6: *Strong complementarity* \Rightarrow *complementarity*.

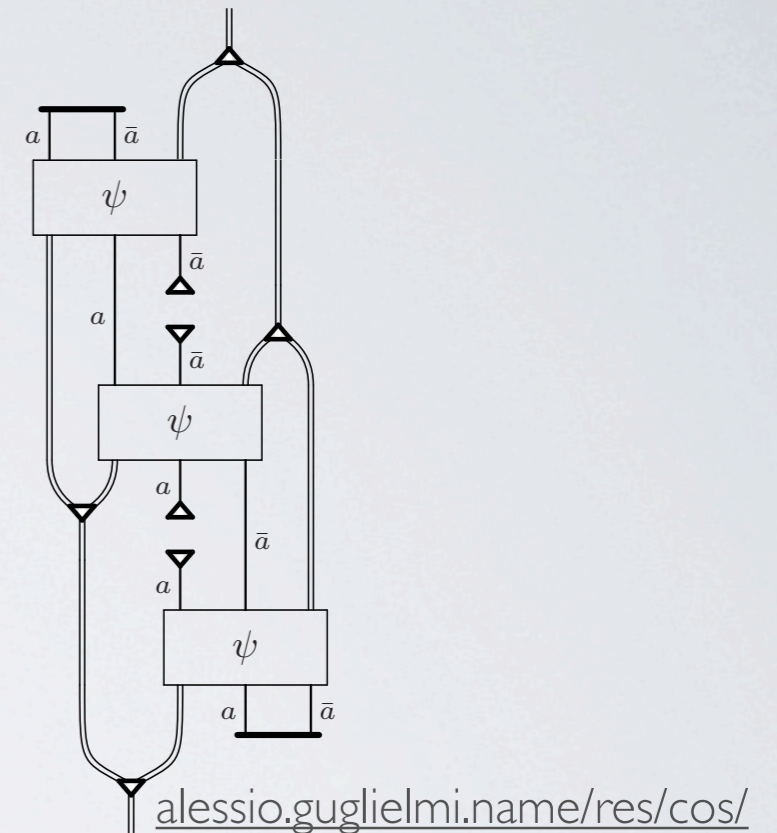
Proof:



arxiv.org/abs/1203.4988

USES, RESEARCH

- More things (off the top of my head):
 - Girard's *Proof nets*
 - Guglielmi's *Atomic flows for deep inference*
 - Lafont's *Algebraic theory of boolean circuits*



| | | | |
|---|---|---|---|
| 1 | 0 | 1 | 1 |
| 1 | 0 | 0 | 0 |
| 1 | 1 | 0 | 1 |

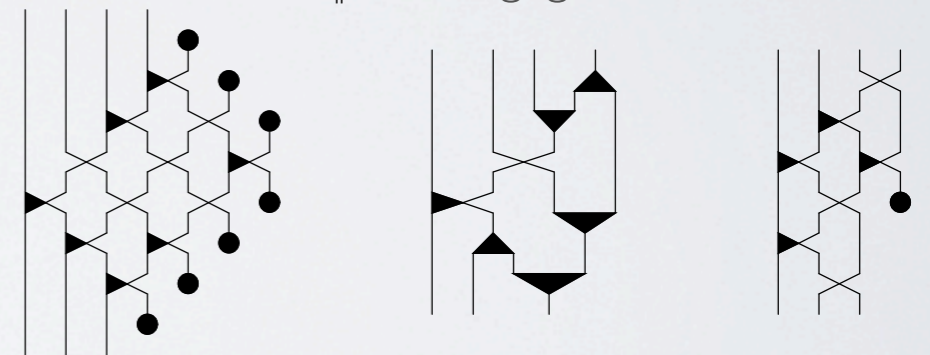


Figure 23: The canonical forms of a matrix in $\mathbf{L}(\mathbb{Z}_2)$

iml.univ-mrs.fr/~lafont/pub/circuits.pdf

Want to draw nice string diagrams for LaTeX? Check out Aleks Kissinger's cross-platform GUI front-end to TikZ, *TikZiT*: tikzit.sourceforge.net/